

# On the Finsler and Doner-Tarski Arithmetical Hierarchies.

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## On the Finsler and Doner-Tarski Arithmetical Hierarchies<sup>1)</sup>

by HILBERT LEVITZ (New York University)

In [5], [1; pg. 64] FINSLER set up a transfinite sequence of binary operations  $\{h_c\}$ <sup>2)</sup> on the ordinals. The first four operations are  $h_0(a, b) = a + 1$ ,  $h_1(a, b) = a + b$ ,  $h_2(a, b) = a \cdot b$ , and  $h_3(a, b) = a^b$ . The hierarchy satisfies the recursion formula  $h_{c+1}(a, b+1) = h_c(h_{c+1}(a, b), a)$ ; this generalizes the formulae  $a + (b+1) = (a+b) + 1$ ,  $a \cdot (b+1) = a \cdot b + a$ ,  $a^{b+1} = a^b \cdot a$ . DONER and TARSKI [3] have also set up a hierarchy  $\{g_c\}$ <sup>3)</sup> of binary operations where each operation is related to the succeeding operation by this same recursion formula, and which, moreover, has a rather simple definition:  $g_0(a, b) = a + b$ ,  $g_c(a, b) = \bigcup_{d < b, e < c} [g_e(g_c(a, d), a)]$  for  $c > 0$ . For every  $c$  and every  $a \neq 0$ ,  $g_c(a, x)$  is a continuous strictly increasing function of  $x$ . On the other hand, this can fail to be the case for  $h_c(a, x)$  if  $c$  is a limit ordinal. Our basic result is that the Doner-Tarski hierarchy is *essentially* what one would get if one deleted from Finsler's hierarchy those operations whose subscript is a limit ordinal. We show also that this deletion does not diminish the representation power of the hierarchy in the sense that if  $d = h_c(a, b)$  for some  $a, b, c, < d$ , then  $d = g_{c'}(a', b')$  for some  $a', b', c', < d$ .

Let  $\{f_c\}$  be the hierarchy of continuous increasing functions defined inductively by:  $f_0(x) = \omega^x$ , if  $c \neq 0$   $f_c$  enumerates in order those ordinals which are fixed points of  $f_s$  for all  $s < c$  (existence proof given in [1]). The fixed points of the function  $f_x(0)$  are called the *strongly critical epsilon numbers*. In [6] we showed that the least strongly critical epsilon number  $\kappa_0$  is the least number greater than  $\omega$  which is inaccessible by means of Finsler's hierarchy; by that we mean the least number  $d > \omega$  such that  $a, b, c, < d$  implies  $h_c(a, b) < d$ . Using the results announced in the above paragraph we will show that  $\kappa_0$  plays the same role with respect to the Doner-Tarski hierarchy. FEFERMAN [4], SCHÜTTE [9], [10], and TAIT [11] have obtained results which show that  $\kappa_0$  plays a significant role in ramified type theory.

**DEFINITION.** To each ordinal  $c$  we associate an ordinal  $c^*$  as follows:  $c^* = c + 1$  if  $c = d + n$  where  $d$  is a limit ordinal and  $0 \leq n < \omega$ ,  $c^* = c$  otherwise. It is easy to see that:

$$c^* + 1 = (c + 1)^* \tag{1}$$

**THEOREM 1.** If  $c \geq 4$  and  $a \geq \omega$  then  $g_{-1+c}(a, 1+b) = h_{c^*}(a, b)$ .

<sup>1)</sup> This work was supported by a grant from the Office of Scientific Research of the United States Air Force.

<sup>2)</sup> We write  $h_c(a, b)$  for Finsler's  $\phi_c(b, a)$  (Note the interchange of the variables). Finsler's restriction that  $a, b, c < \Omega_1$  can be lifted [1; pg. 64].

<sup>3)</sup> We write  $g_c(a, b)$  for Doner-Tarski's  $a0_c b$ .

*Proof:* By transfinite induction on  $\Omega_i c + b$ . Our induction hypothesis is that  $c \geq 4$  and  $\Omega_i c' + b' < \Omega_i c + b$  implies that  $g_{-1+c'}(a, 1+b') = h_{(c')^*}(a, b')$  for all  $a \geq \omega$ .<sup>4)</sup> We must show under this assumption that  $g_{-1+c}(a, 1+b) = h_{c^*}(a, b)$  for all  $a \geq \omega$ . Let  $a \geq \omega$  be given:

CASE 1.  $c=4$ ; using [3; 3(iii)] and [5; pg. 80] we get:

$$g_{-1+c}(a, 1+b) = g_3(a, 1+b) = a^{(a^b)} = h_4(a, b) = h_c(a, b) = h_{c^*}(a, b).$$

CASE 2.  $c > 4$ ;

CASE 2.1  $b=0$ ; using [3; 17(iii)] and [5; th. 9] we get:

$$g_{-1+c}(a, 1+b) = g_{-1+c}(a, 1) = a = h_{c^*}(a, 0) = h_{c^*}(a, b).$$

CASE 2.2  $b \neq 0$ ;

CASE 2.2.1  $b$  is a limit ordinal; using [3; 17(vi)], our induction hyp. and the fact that  $h_{c^*}(a, x)$  is a continuous increasing function of  $x$  [1; pg. 65] we get:

$$g_{-1+c}(a, 1+b) = \sup_{d < b} g_{-1+c}(a, 1+d) = \sup_{d < b} h_{c^*}(a, d) = h_{c^*}(a, b).$$

CASE 2.2.2  $b$  is a successor ordinal  $v+1$ ;

CASE 2.2.2.1  $c$  is a successor ordinal  $d+1$ ; then using [3; 17(iv)] and the fact that  $a \geq \omega$  we get:

$$\begin{aligned} g_{-1+c}(a, 1+b) &= g_{-1+(d+1)}(a, 1+(v+1)) = g_{(-1+d)+1}(a, (1+v)+1) \\ &= g_{-1+d}(g_{(-1+d)+1}(a, 1+v), a) = g_{-1+d}(g_{(-1+d)+1}(a, 1+v), 1+a). \end{aligned}$$

Now by [3; corr. 5(i)] we note that  $g_{-1+d}(a, 1+v) \geq a \geq \omega$ , so using successive applications of our induction hypothesis followed by (1), [5; pg. 6] and (1) again, we get:

$$\begin{aligned} g_{-1+d}(g_{(-1+d)+1}(a, 1+v), 1+a) &= h_{d^*}(g_{-1+(d+1)}(a, 1+v), a) \\ &= h_{d^*}(h_{(d+1)^*}(a, v), a) = h_{d^*}(h_{d^*+1}(a, v), a) \\ &= h_{d^*+1}(a, v+1) = h_{(d+1)^*}(a, b) = h_{c^*}(a, b). \end{aligned}$$

CASE 2.2.2.2.  $c$  is a limit ordinal; using [3; 17(v)] we get:

$$\begin{aligned} g_{-1+c}(a, 1+b) &= g_c(a, 1+(v+1)) = g_c(a, (1+v)+1) \\ &= \sup_{d < c} g_d(g_c(a, 1+v), a) = \sup_{d < c} g_{-1+d}(g_c(a, 1+v), a) \\ &= \sup_{d < c} g_{-1+d}(g_{-1+c}(a, 1+v), a). \end{aligned}$$

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<sup>4)</sup> Our induction is up to  $\Omega_i^2$ , where  $\Omega_i$  is an initial ordinal and  $a, b, c < \Omega_i$ .

Now by [3; corr. 5(i)] and induction hyp.  $\omega \leq g_{-1+c}(a, 1+v) = h_{c^*}(a, v)$  so

$$\sup_{d < c} g_{-1+d}(g_{-1+c}(a, 1+v), a) = \sup_{d < c} g_{-1+d}(h_{c^*}(a, v), a) = \sup_{d < c} g_{-1+d}(h_{c^*}(a, v), 1+a).$$

Now by a further application of the induction hypothesis, Finsler’s definition, and [5, pg. 6] we get:

$$\begin{aligned} \sup_{d < c} g_{-1+d}(h_{c^*}(a, v), 1+a) &= \sup_{d < c} h_{d^*}(h_{c^*}(a, v), a) \\ &= \sup_{d < c} h_d(h_{c^*}(a, v), a) = h_c(h_{c^*}(a, v), a) = h_c(h_{c+1}(a, v), a) \\ &= h_{c+1}(a, v+1) = h_{c^*}(a, b) \quad (\text{Q.E.D.}). \end{aligned}$$

**COROLLARY 1.** *If  $a, b, c, < \kappa_0$  then  $g_c(a, b) < \kappa_0$ .*

*Proof:* Follows from the main theorem together with the fact that Finsler’s hierarchy has the same property [6, th. 3].

**COROLLARY 2.** *If  $x \geq 1$ , then*

$$f_c(-1+x) = \begin{cases} g_{2c+2}(\omega, \omega x) & \text{if } 1 \leq c < \omega \\ g_c(\omega, 1+x) & \text{if } c \text{ is a limit ordinal} \\ g_{a+2n}(\omega, \omega x) & \text{if } c = a+n \text{ where } a \text{ is a limit} \\ & \text{ordinal and } 1 \leq n < \omega \end{cases}$$

*Proof:* Follows from the main theorem together with [6, th. 1].

**THEOREM 2.** *If  $d > \omega$  and  $d = h_c(a, b)$  where  $a, b, c, < d$ , then  $d = g_u(v, w)$  where  $u, v, w < d$ .*

*Proof:* By [6, th. 3]  $d \neq f_a(0)$ ; but then by [8; (5.1)]  $d = a+b$  where  $a, b < d$  or  $d = f_c(a)$  where  $c, a < d$ :

**CASE 1.**  $d = a+b$  where  $a, b < d$ ; then  $d = g_0(a, b)$ .

**CASE 2.**  $d = f_c(a)$  where  $c, a < d$ ; write  $a = -1+b$  where  $b \geq 1$ :

**CASE 2.1**  $c \neq 0$ ; then  $d$  is an epsilon number, so  $a < d$  implies  $b < d, 1+b < d$ , and  $\omega b < d$ .

**CASE 2.1.1**  $c$  is a limit ordinal; then by corollary 2:

$$c, 1+b, \omega < d = f_c(-1+b) = g_c(\omega, 1+b).$$

**CASE 2.1.2**  $c = e+n$  where  $e$  is a limit ordinal and  $1 \leq n < \omega$ ; then by corollary 2:

$$e+2n, \omega, \omega b < d = f_c(-1+b) = g_{e+2n}(\omega, \omega b).$$

**CASE 2.1.3**  $c$  is finite; then by corollary 2:

$$2c+2, \omega, \omega b < d = f_c(-1+b) = g_{2c+2}(\omega, \omega b).$$

CASE 2.2  $c=0$ ; now  $a \geq 2$  because  $\omega < d = f_0(a) = \omega^a$ , thus by [3, 3(ii)]  $d = f_2(\omega, a) > 2, a, \omega$ . (Q.E.D.)

COROLLARY 3.  $\kappa_0$  (the least strongly critical epsilon number) is the least ordinal inaccessible by means of the Doner–Tarski hierarchy.

*Proof:*  $\kappa_0$  is inaccessible by corollary 1. That it is the least such number follows from the main theorem together with the result of [6] that  $\kappa_0$  is the least ordinal greater than  $\omega$  which is inaccessible by means of Finsler's hierarchy.

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