

2. The Division Algorithm

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If R is a field, then $w: \text{deg}(J) \rightarrow R$ is a "weight function".

$$\delta \mapsto 1$$

So the corresponding figure is of the form

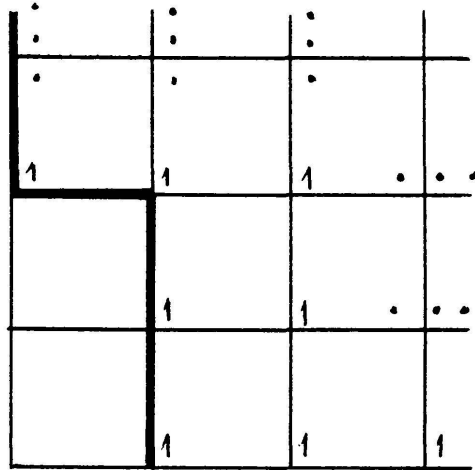


FIGURE 3.

2. THE DIVISION ALGORITHM

Let F be a finite subset of $R[X] - \{0\}$.

2.1. *Definition.* An "admissible combination of F " is an expression of the form $L := \sum_{\gamma \in \mathbb{N}^n, P \in F} c(\gamma, P)X^\gamma P$, $c(\gamma, P) \in R$, such that

$$\text{deg}(L) = \max \{ \text{deg}(X^\gamma P) \mid c(\gamma, P) \neq 0 \}.$$

Example. Let $P, Q \in R[X]$ and let $\alpha, \beta \in \mathbb{N}^n$. Then $X^\alpha P - X^\beta Q$ is an admissible combination of $\{P, Q\}$ iff $X^\alpha \cdot \text{in}(P) \neq X^\beta \cdot \text{in}(Q)$.

Remark. For every $Q \in \langle \text{in}(F) \rangle$ there is an admissible combination L of F such that $\text{in}(L) = \text{in}(Q)$. L can be calculated in the following way:

Let $F' := \{P \in F \mid \text{deg}(Q) - \text{deg}(P) \in \mathbb{N}^n\}$. Then

$$Q \in \langle \text{in}(F') \rangle \quad \text{and} \quad \text{lc}(Q) \in_R \langle \text{lc}(P) \mid P \in F' \rangle.$$

For $P \in F'$ we calculate elements $c(P) \in R$ such that $\text{lc}(Q) = \sum_{P \in F'} c(P) \text{lc}(P)$.

Set $L := \sum_{P \in F'} c(P)X^{\text{deg}(Q) - \text{deg}(P)}P$.

Example: $F := \{5X_1 + 1, 3X_2 + 2\}$, $Q := X_1^2 X_2^3$.

Then $L = -X_1 X_2^3 (5X_1 + 1) + 2X_1^2 X_2^2 (3X_2 + 2)$.

2.2. PROPOSITION. Every $Q \in R[X] - \{0\}$ may be written as $Q = L + \bar{Q}$ with the following properties:

If $\text{in}(Q) \notin \langle \text{in}(F) \rangle$, then $L = 0$ and $Q = \bar{Q}$.

If $\text{in}(Q) \in \langle \text{in}(F) \rangle$, then L is an admissible combination of F with $\text{in}(L) = \text{in}(Q)$, and either $\bar{Q} = 0$ or $\text{in}(\bar{Q}) \notin \langle \text{in}(F) \rangle$.

L and \bar{Q} can be found in a finite number of steps by the following algorithm:

$$Q_0 := Q;$$

For $k \in \mathbf{N}$ assume that Q_k has already been defined. If $\text{in}(Q_k) \in \langle \text{in}(F) \rangle$, we define $Q_{k+1} := Q_k - L_k$, where L_k is an admissible combination of F with $\text{in}(L_k) = \text{in}(Q_k)$.

If $Q_k = 0$ or $\text{in}(Q_k) \notin \langle \text{in}(F) \rangle$, then $L := \sum_{j=0}^{k-1} L_j$ and $\bar{Q} := Q_k$.

Proof. We only have to show that there is a number $k \in \mathbf{N}$ such that $\text{in}(Q_k) \notin \langle \text{in}(F) \rangle$ or $Q_k = 0$.

If $\text{in}(Q_j) \in \langle \text{in}(F) \rangle$, then $\deg(Q_j) > \deg(Q_{j+1})$, so the assertion follows from the lemma 1.3.

2.3. Definition. The algorithm above is called "division by F ", the polynomial \bar{Q} (or, more precisely, \bar{Q}^F) is "a rest of Q after division by F ".

Remarks.

- 1) Even if the strict ordering $<$ is fixed, \bar{Q} depends on the choice of the L_k 's in the algorithm. Hence \bar{Q} is in general not uniquely determined by Q and F .
- 2) If a rest of Q after division by F is zero, then Q belongs to the ideal generated by F . In general the inverse is not true.

2.4. Example. Consider the graded lexicographic ordering and

$$P_1 := 2X_1^2 + X_1X_2, \quad P_2 := 3X_2^2 + X_1 \in \mathbf{Z}[X_1, X_2].$$

Let F be $\{P_1, P_2\}$ and let $Q := 2X_1^3X_2^3 + X_1X_2$. Then $Q_0 = Q$.

$$L_0 := 2X_1^3X_2P_2 - 2X_1X_2^3P_1,$$

$$Q_1 := Q_0 - L_0 = -2X_1^2X_2^4 + 2X_1^4X_2 + X_1X_2.$$

$$L_1 := -2X_1^2X_2^2P_2 + 2X_2^4P_1,$$

$$Q_2 := Q_1 - L_1 = -2X_1X_2^5 + 2X_1^4X_2 + 2X_1^3X_2^2 + X_1X_2.$$

Now in $(Q_2) \notin \langle \text{in}(F) \rangle$, therefore $Q = L_0 + L_1 + Q_2$.

See figure 4.

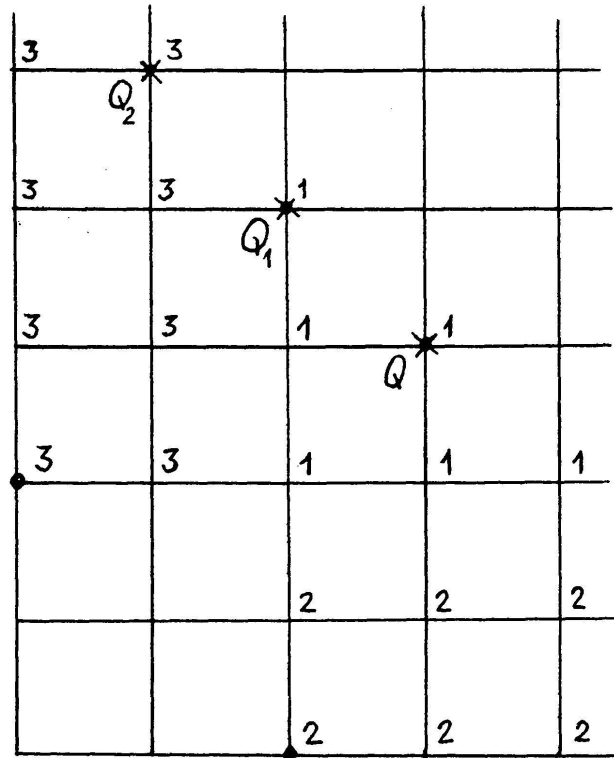


FIGURE 4.

But if we choose $L'_0 := X_1 X_2^3 P_1$, then

$$Q'_1 := Q_0 - L'_0 = -X_1^2 X_2^4 + X_1 X_2,$$

$$L'_1 := -X_1^2 X_2^2 P_2 + X_2^4 P_1$$

$$Q'_2 := Q'_1 - L'_1 = -X_1 X_2^5 + X_1^3 X_2^2 + X_1 X_2,$$

$$\text{therefore } Q = L'_0 + L'_1 + Q'_2.$$

So Q_2 and Q'_2 are rests of Q after division by F and $Q_2 \neq Q'_2$.

2.5. PROPOSITION. Let J be an ideal in $R[X]$ containing F . Then the following conditions are equivalent:

- (1) F is a Gröbner basis of J .
- (2) For every $Q \in J$, each rest of Q after division by F is zero.
- (3) For every $Q \in J$, a rest of Q after division by F is zero.

Proof.

(1) \Rightarrow (2): Division of $Q \in J$ by F yields $Q = L + \bar{Q}$ with $\bar{Q} = 0$ or $\text{in}(\bar{Q}) \notin \langle \text{in}(F) \rangle$. Now $L \in J$ and $Q \in J$ imply $\bar{Q} \in J$. Since $\langle \text{in}(J) \rangle = \langle \text{in}(F) \rangle$, \bar{Q} must be zero.

(2) \Rightarrow (3): trivial.

(3) \Rightarrow (1): By (3) we have $\text{in}(Q) \in \langle \text{in}(F) \rangle$ for every $Q \in J - \{0\}$. Hence $\langle \text{in}(J) \rangle = \langle \text{in}(F) \rangle$.

2.6. COROLLARY. Let F be a Gröbner basis of an ideal $J \leq R[X]$.

1) F generates J .

2) Let $Q \in R[X]$. Then $Q \in J$ iff a rest of Q after dividing by F is zero.

Proof. Obvious.

2.7. Another characterisation of Gröbner bases can be given as follows:

We shall say that a set $\{L_\alpha \mid \alpha \in \mathcal{D}(F)\}$ of admissible combinations of F (with pairwise different degrees) is an " F -admissible set", if for all α we have $\deg(L_\alpha) = \alpha$ and $\text{lc}(L_\alpha)$ generates the ideal

$${}_R \langle \text{lc}(P) \mid P \in \langle \text{in}(F) \rangle, \deg(P) = \alpha \rangle .$$

Any F -admissible set is R -linearly independent.

If R is a field the condition on $\text{lc}(L_\alpha)$ is superfluous.

PROPOSITION. Let J be an ideal in $R[X]$ containing F . Then the following conditions are equivalent:

- (1) F is a Gröbner basis of J .
- (2) There is an F -admissible set which is a R -basis of J .
- (3) Every F -admissible set is a R -basis of J .

Proof. Let $\{L_\alpha \mid \alpha \in \mathcal{D}(F)\}$ be a F -admissible set.

(1) \Rightarrow (3): Let Q be an element of $J - \{0\}$. Division of Q by $\{L_{\deg(Q)}\}$, of its rest \bar{Q} by $\{L_{\deg(\bar{Q})}\}$, ... yields in a finite number of steps an expression of Q as R -linear combination of L_α 's.

(3) \Rightarrow (2): trivial.

(2) \Rightarrow (1): Suppose that $\{L_\alpha \mid \alpha \in \mathcal{D}(F)\}$ is a R -basis of J . For every $Q \in J - \{0\}$ the initial term of $L_{\deg(Q)}$ divides $\text{in}(Q)$, hence $\text{in}(Q) \in \langle \text{in}(F) \rangle$.

3. CONSTRUCTION OF GRÖBNER BASES

3.1. *Definition.* Let P, Q be elements of $R[X]$, let $\alpha, \beta \in \mathbb{N}^n$ and let $a, b \in R$. Then the polynomial